Geometric data structures

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Introduction to Graph and Geometric Algorithms

Scope of the lecture

- Binary search trees and 2-d range trees
  We consider 1-d and 2-d range queries for point sets.

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  Interval trees for reporting all (horizontal) intervals containing a given (vertical) query line or segment.

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- **Planar point location**
  Using triangulation refinement and monotone subdivisions.

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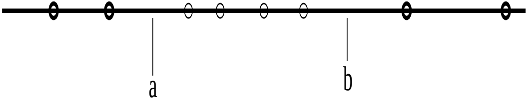
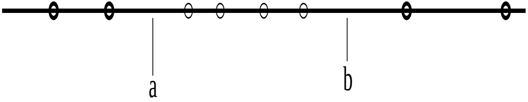
- **Interval trees**
  Interval trees for reporting all (horizontal) intervals containing a given (vertical) query line or segment.
- **Planar point location**
  Using triangulation refinement and monotone subdivisions.
- **Hierarchical representation of a convex polygon**
  Detecting the intersection of a convex polygon with a query line..

1-dimensional Range searching



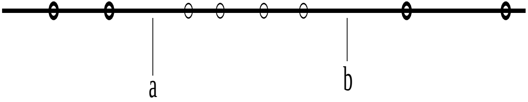
- Problem: Given a set P of n points $\{p_1, p_2, \cdots, p_n\}$ on the real line, report points of P that lie in the range $[a, b]$, $a \leq b$.

1-dimensional Range searching



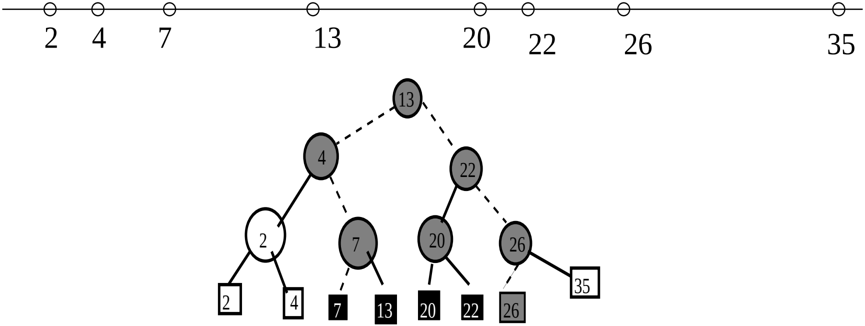
- Problem: Given a set P of n points $\{p_1, p_2, \cdots, p_n\}$ on the real line, report points of P that lie in the range $[a, b]$, $a \leq b$.
- Using binary search on an array we can answer such a query in $O(\log n + k)$ time where k is the number of points of P in $[a, b]$.

1-dimensional Range searching



- Problem: Given a set P of n points $\{p_1, p_2, \cdots, p_n\}$ on the real line, report points of P that lie in the range $[a, b]$, $a \leq b$.
- Using binary search on an array we can answer such a query in $O(\log n + k)$ time where k is the number of points of P in $[a, b]$.
- However, when we permit insertion or deletion of points, we cannot use an array answering queries so efficiently.

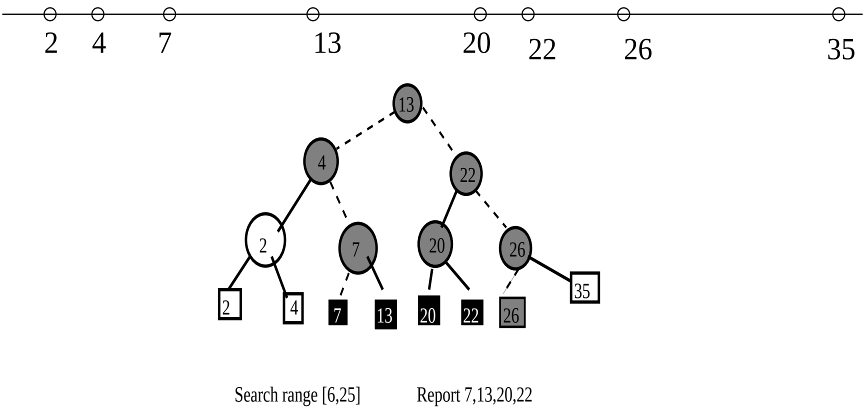
1-dimensional Range searching



Search range [6,25] Report 7,13,20,22

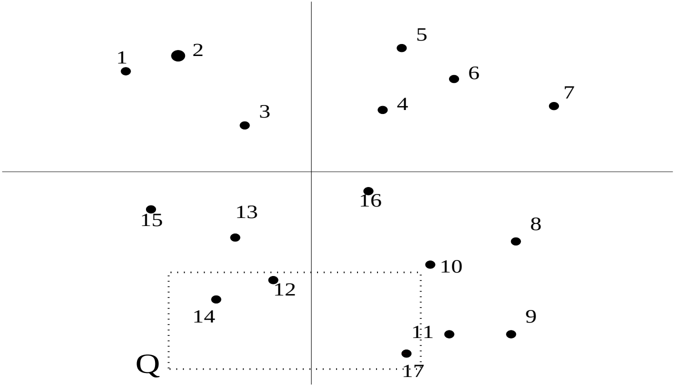
- We use a *binary leaf search tree* where leaf nodes store the points on the line, sorted by x-coordinates.

1-dimensional Range searching



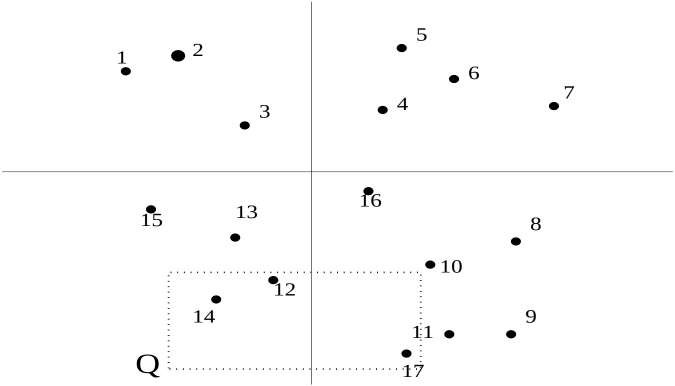
- We use a *binary leaf search tree* where leaf nodes store the points on the line, sorted by x-coordinates.
- Each internal node stores the x-coordinate of the rightmost point in its left subtree for guiding search.

2-dimensional Range Searching



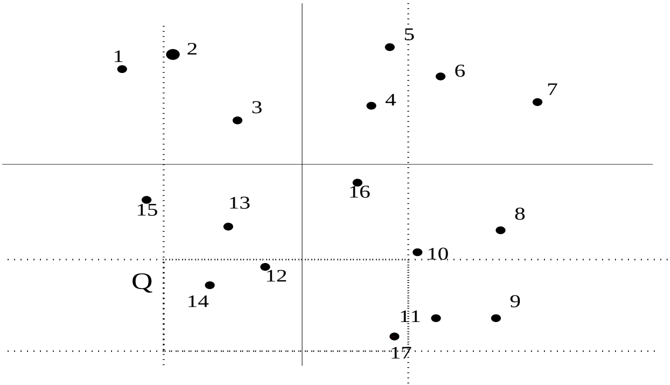
- Problem: Given a set P of n points in the plane, report points inside a query rectangle Q whose sides are parallel to the axes.

2-dimensional Range Searching



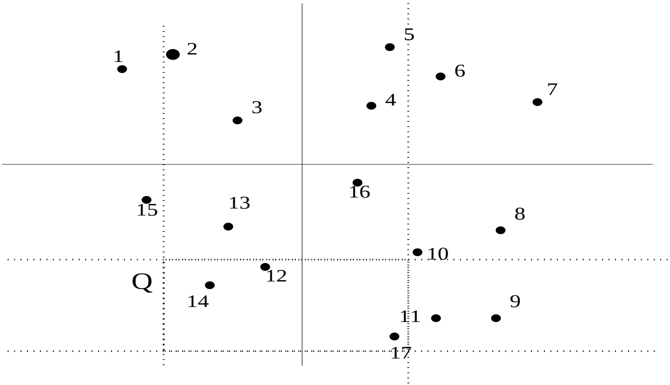
- Problem: Given a set P of n points in the plane, report points inside a query rectangle Q whose sides are parallel to the axes.
- Here, the points inside R are 14, 12 and 17.

2-dimensional Range Searching



- Using two 1-d range queries, one along each axis, solves the 2-d range query.

2-DIMENSIONAL RANGE SEARCHING



- Using two 1-d range queries, one along each axis, solves the 2-d range query.
- The cost incurred may exceed the actual output size of the 2-d range query.

Range searching with range trees and Kd-trees

- Given a set S of n points in the plane, we can construct a *2d-range tree* in $O(n \log n)$ time and space, so that rectangle queries can be executed in $O(\log^2 n + k)$ time.

Range searching with range trees and Kd-trees

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- The query time can be improved to $O(\log n + k)$ using the technique of *fractional cascading*.

Range searching with range trees and Kd-trees

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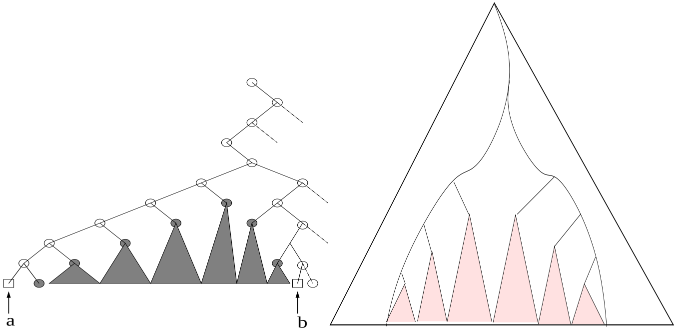
- The query time can be improved to $O(\log n + k)$ using the technique of *fractional cascading*.
- Given a set S of n points in the plane, we can construct a Kd-tree in $O(n \log n)$ time and $O(n)$ space, so that *rectangle queries* can be executed in $O(\sqrt{n} + k)$ time. Here, the number of points in the query rectangle is k.

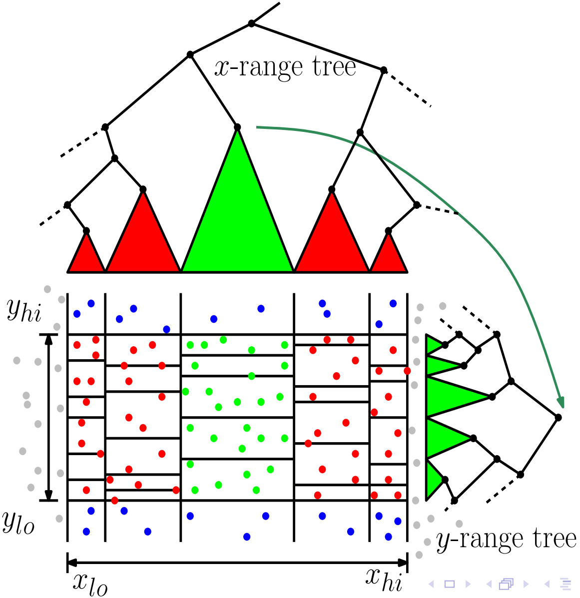
Range searching in the plane using range trees



Given a 2-d rectangle query $[a, b]X[c, d]$, we can identify subtrees whose leaf nodes are in the range $[a, b]$ along the X-direction.

Only a subset of these leaf nodes lie in the range $[c, d]$ along the Y-direction.

Range searching in the plane using range trees



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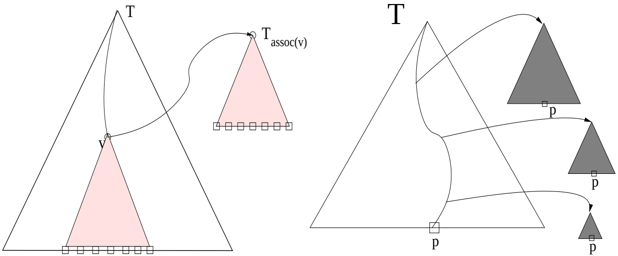
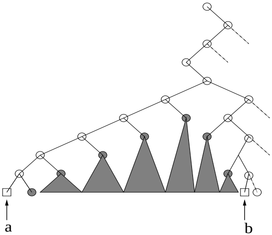


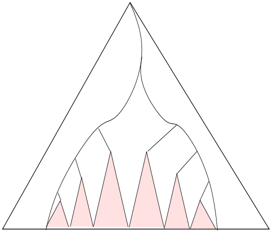
$T_{assoc(v)}$ is a binary search tree on y-coordinates for points in the leaf nodes of the subtree tooted at v in the tree T.

The point p is duplicated in $T_{assoc(v)}$ for each v on the search path for p in tree T.

The total space requirement is therefore $O(n \log n)$.

Range searching in the plane using range trees

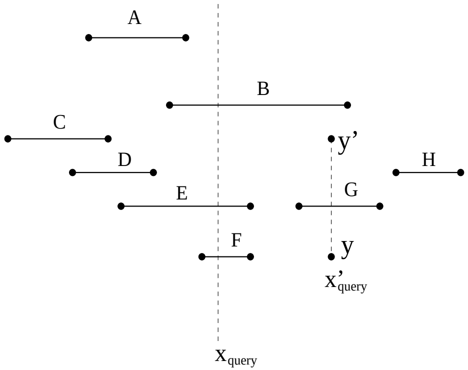




We perform 1-d range queries with the y-range $[c, d]$ in each of the subtrees adjacent to the left and right search paths within the x-range $[a, b]$ in the tree T.

Since the search path is $O(\log n)$ in size, and each y-range query requires $O(\log n)$ time, the total cost of searching is $O(\log^2 n)$. The reporting cost is $O(k)$ where k points lie in the query rectangle.

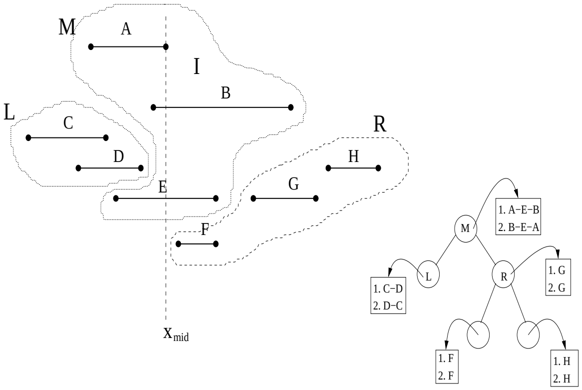
Finding intervals containing a vertical query line/segment



Simpler queries ask for reporting all intervals intersecting the vertical line $X = x_{query}$.

More difficult queries ask for reporting all intervals intersecting a vertical segment joining (x'_{query}, y) and (x'_{query}, y').

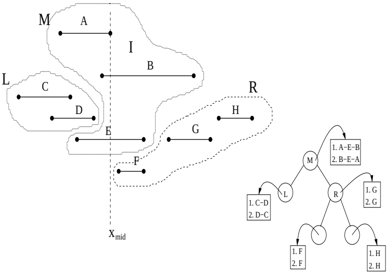
Constructing the interval tree



The set M has intervals intersecting the vertical line $X = x_{mid}$, where x_{mid} is the median of the x-coordinates of the $2n$ endpoints.

The root node has intervals M sorted in two independent orders (i) by right end points (B-E-A), and (ii) left end points (A-E-B).

Answering queries using an interval tree



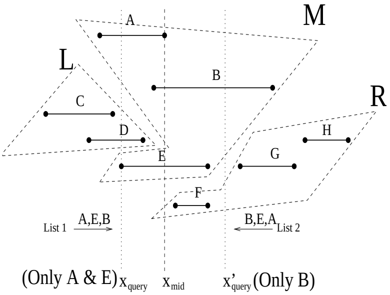
The set L and R have at most n endpoints each.

So they have at most $\frac{n}{2}$ intervals each.

Clearly, the cost of (recursively) building the interval tree is $O(n \log n)$.

The space required is linear.

Answering queries using an interval tree

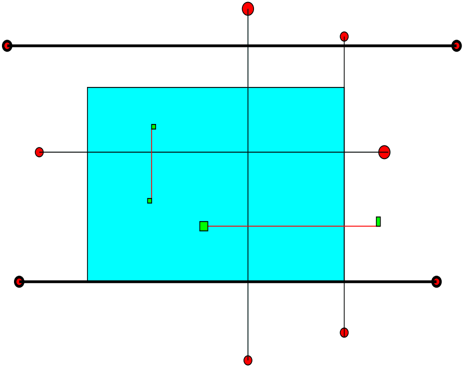


For $x_{query} < x_{mid}$, we do not traverse subtree for subset R.

For $x'_{query} > x_{mid}$, we do not traverse subtree for subset L.

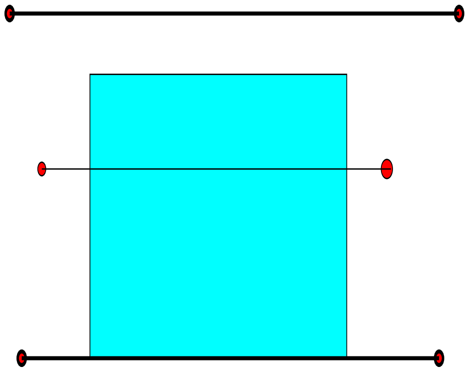
Clearly, the cost of reporting the k intervals is $O(\log n + k)$.

Reporting (portions of) all rectilinear segments inside a query rectangle



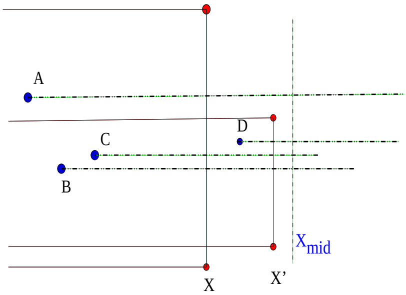
For detecting segments with one (or both) ends inside the rectangle, it is sufficient to maintain rectangular range query apparatus for output-sensitive query processing.

Reporting segments with no endpoints inside the query rectangle



Report all (horizontal) segments that cut across the query rectangle or include an entire (top/bottom) bounding edge.
Use either the right (or left) edge, and the top (or bottom) edge of the query rectangle.

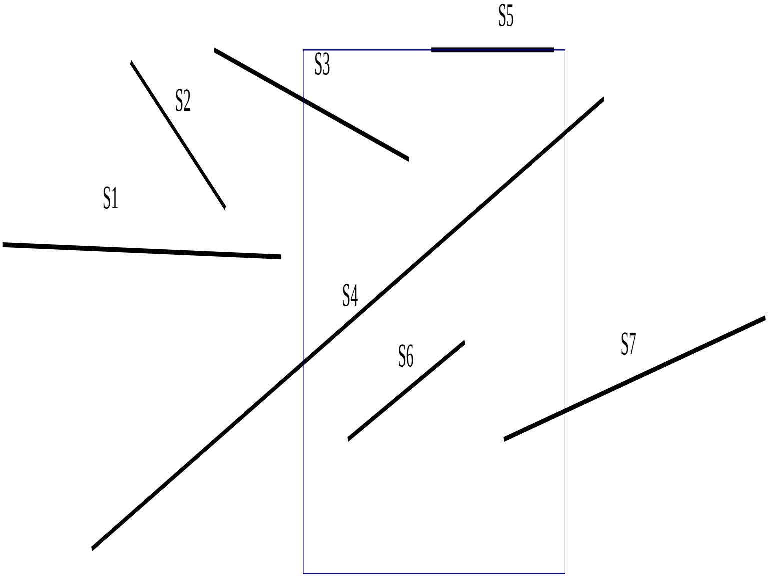
Right edges X and X' of two query rectangles



Use an interval tree of all the horizontal segments and the right bounding edge of the query rectangle like X or X'.

This helps reporting all segments cutting the right edge of the query rectangle.

Use the rectangle query for vertical segment X and find points A, B and C in the rectangle with left edge at minus infinity. For X', report B, C and D, similarly.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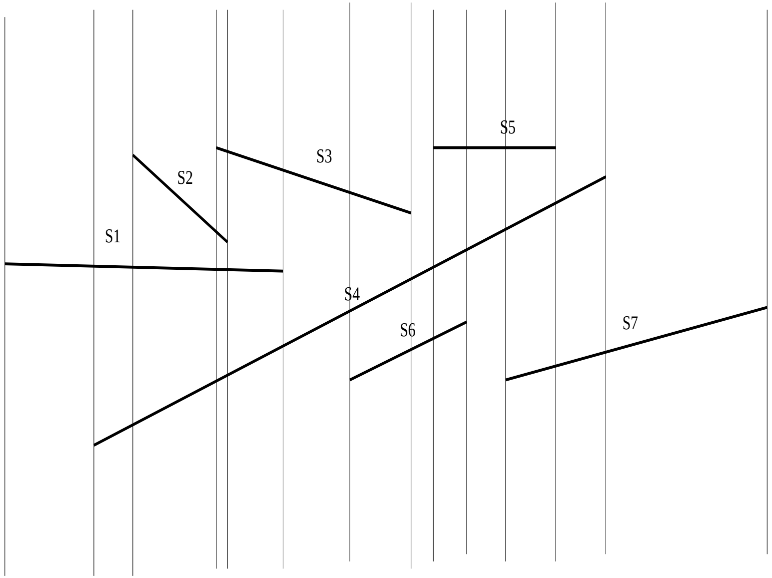
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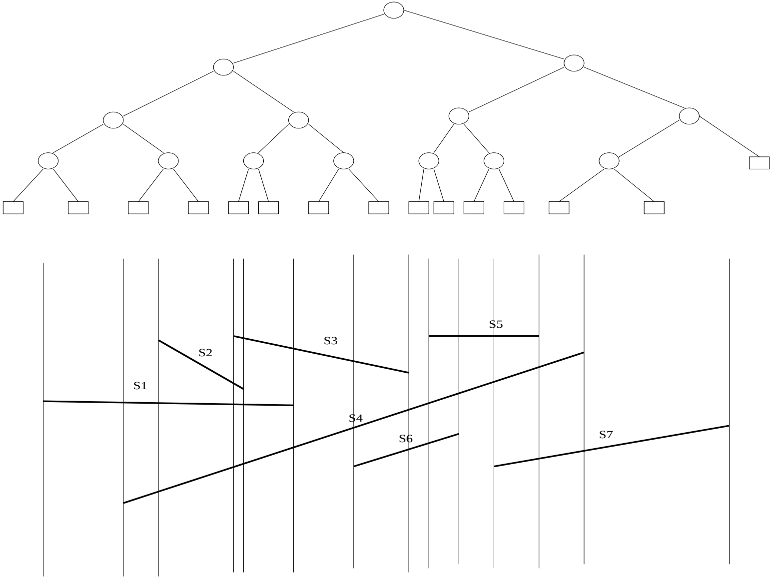
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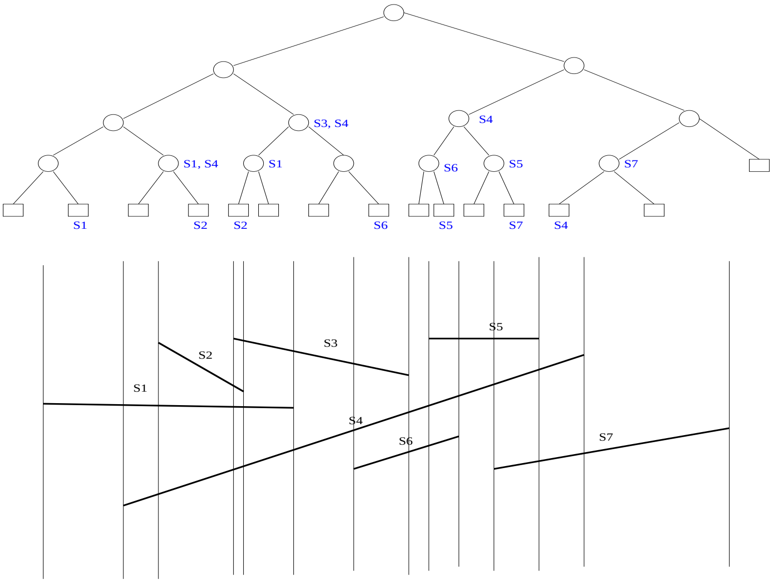

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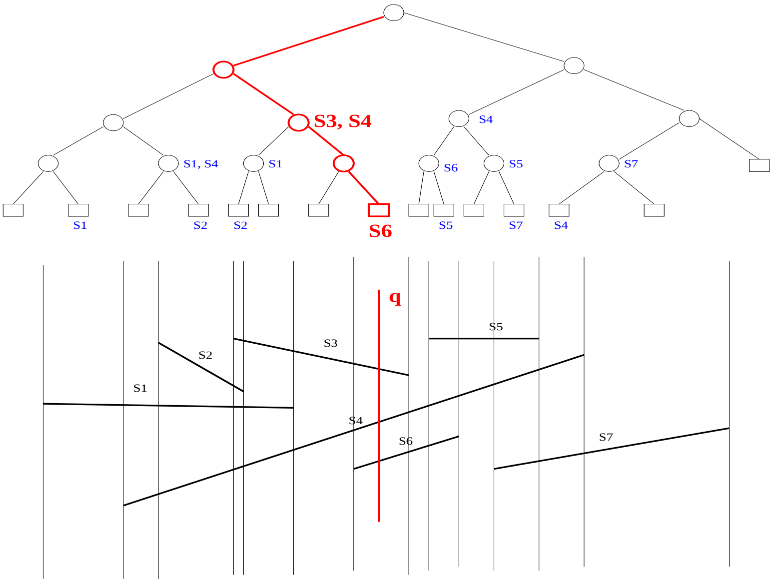
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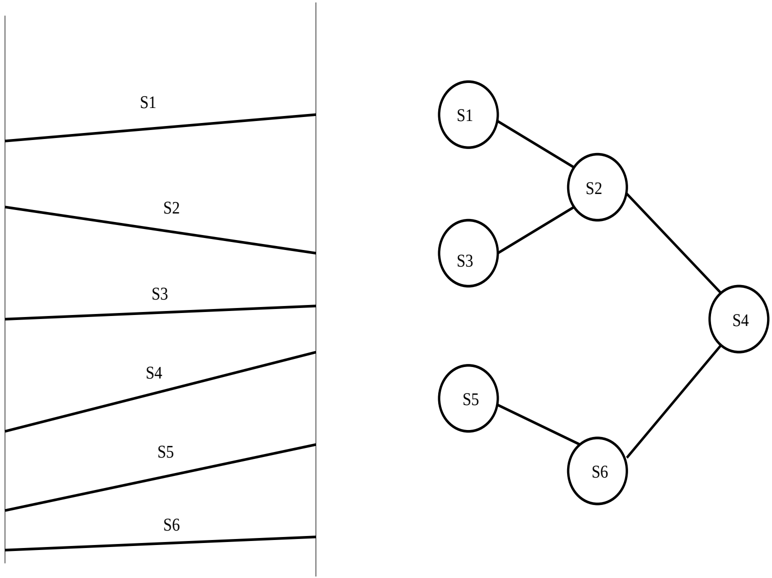
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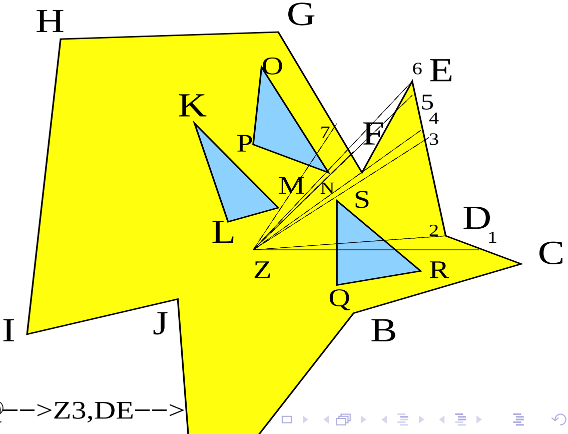
Computing the visible region in a polygon with opaque obstacles

SQ,SR,DC,1−−>SQ,SR,DE,2−−>DE,3−-

FG,FE,DE,4−−>NP,NO,FG,FE,DE,5−−>

NP,NO,FG,FE,DE,6−−>LM,MK,NP,NO,FG,7



Z1 ,SQ−−>Z2,SQ−−>Z3,DE−−>

Computing the visible region in a polygon with opaque obstacles

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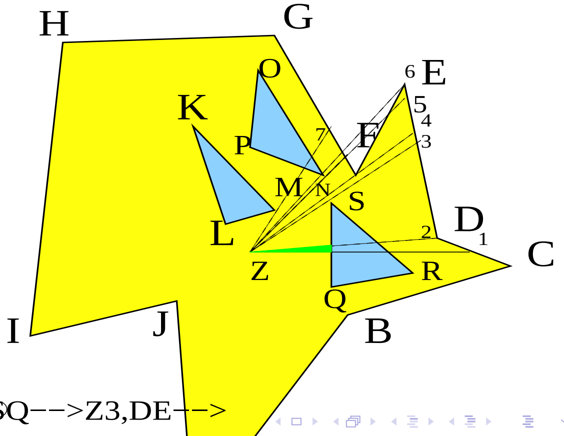

Z1 ,SQ-->Z2,SQ-->Z3,DE-->

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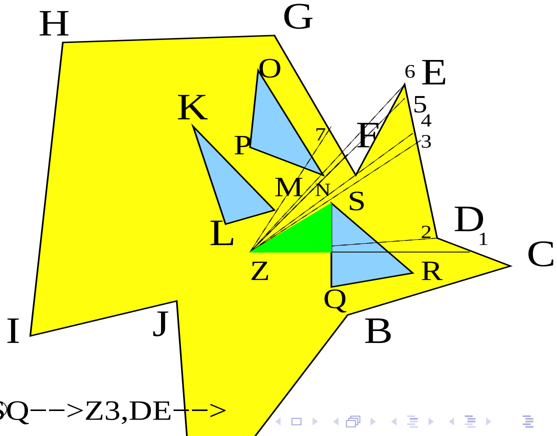

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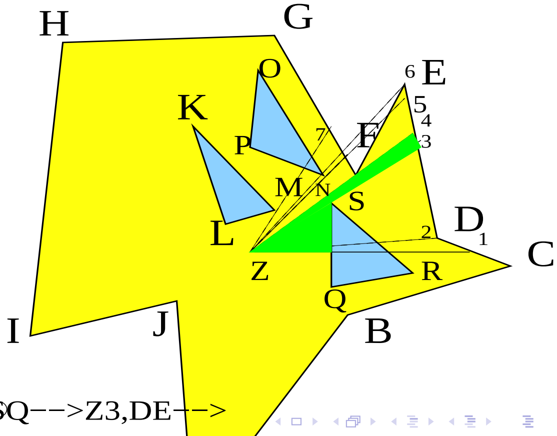

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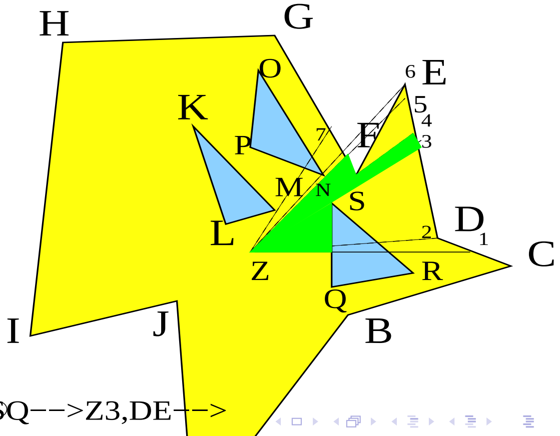

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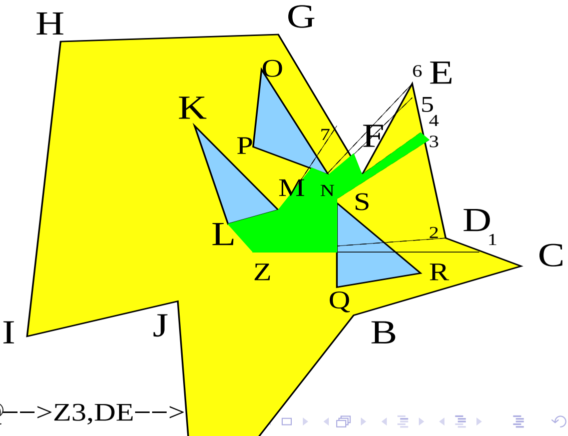

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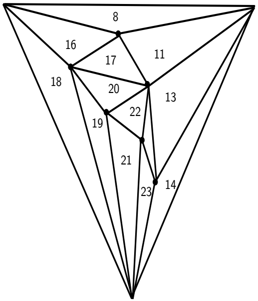
Planar point location by triangulation refinement



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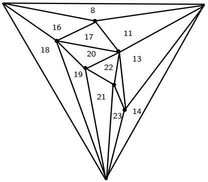

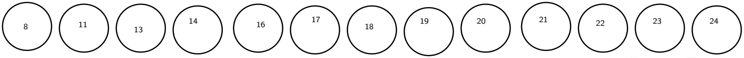

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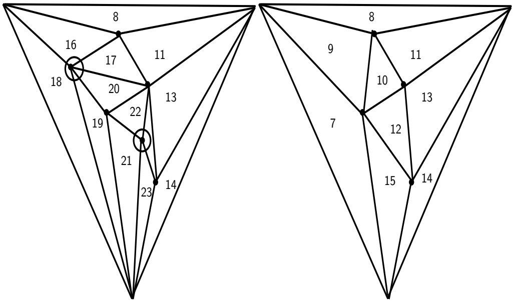

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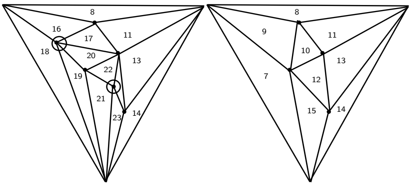

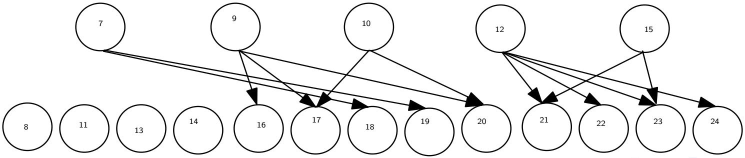

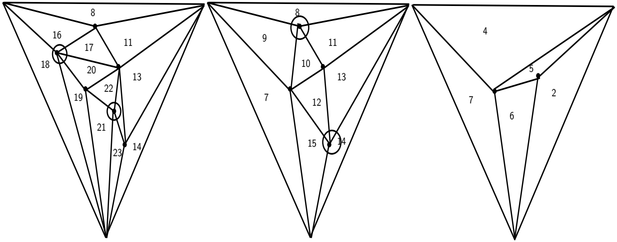
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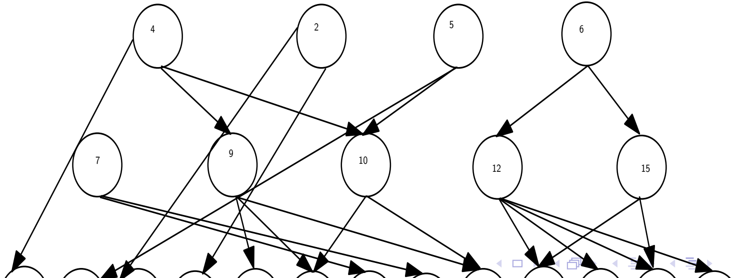
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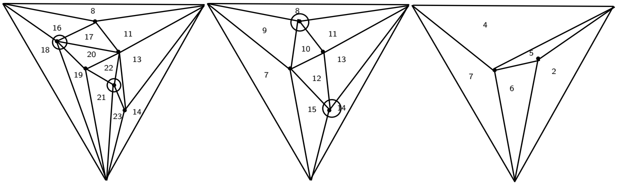
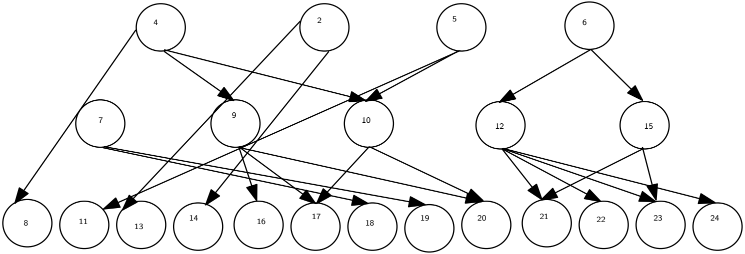
Planar point location by triangulation refinement





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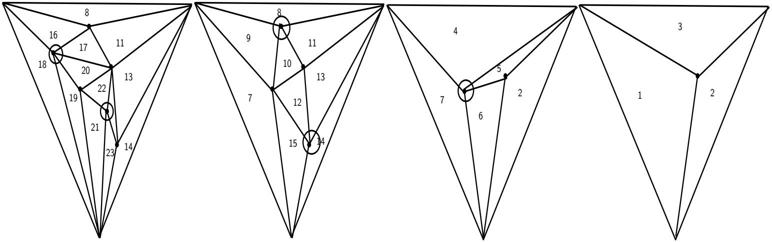

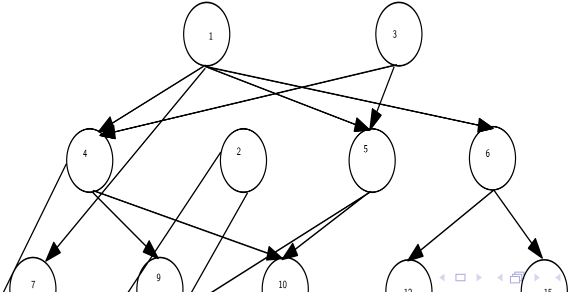

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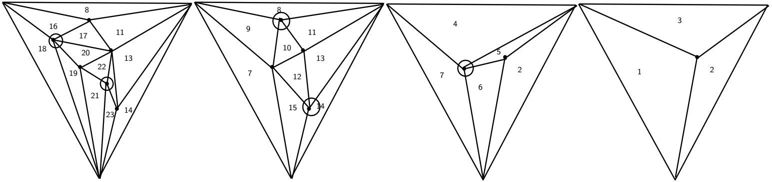
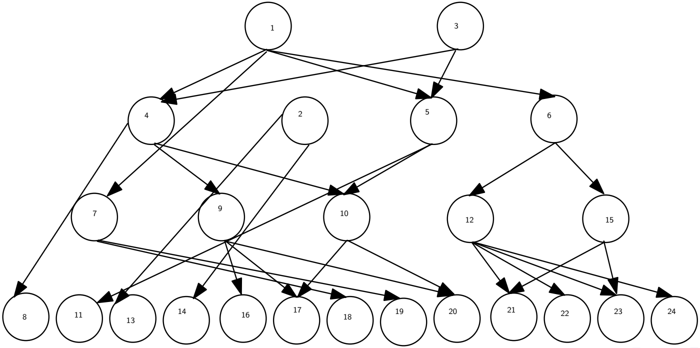
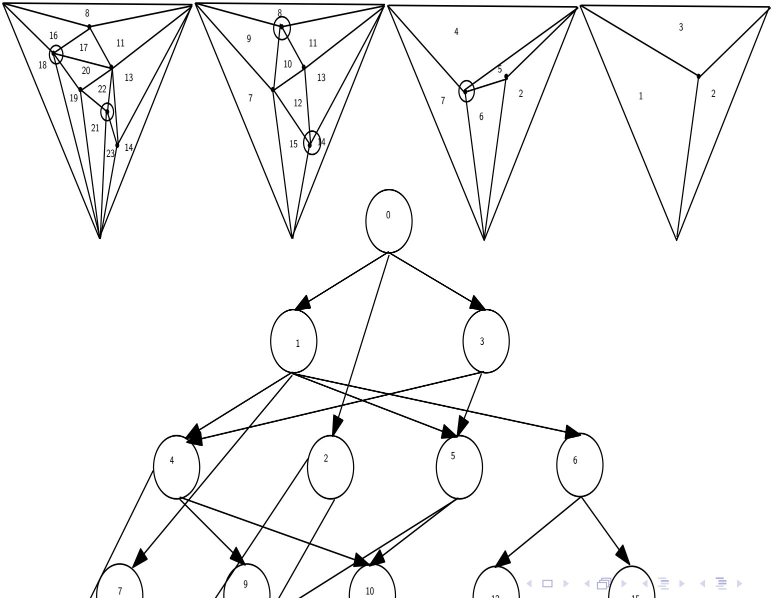
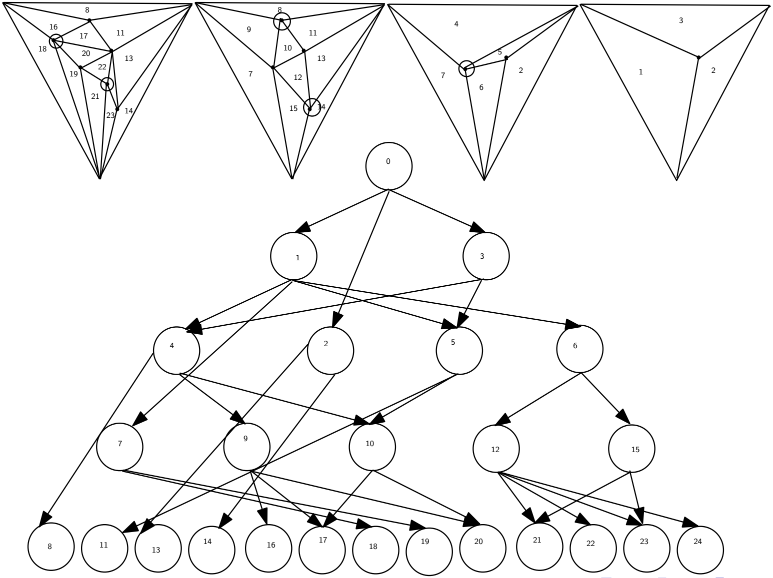




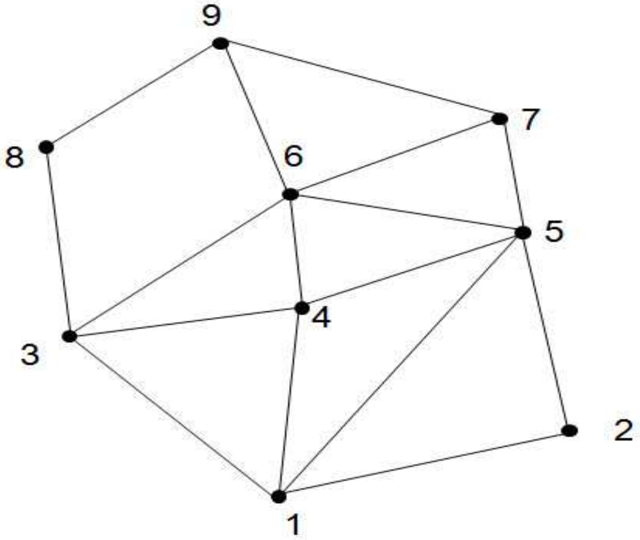
Planar point location by triangulation refinement



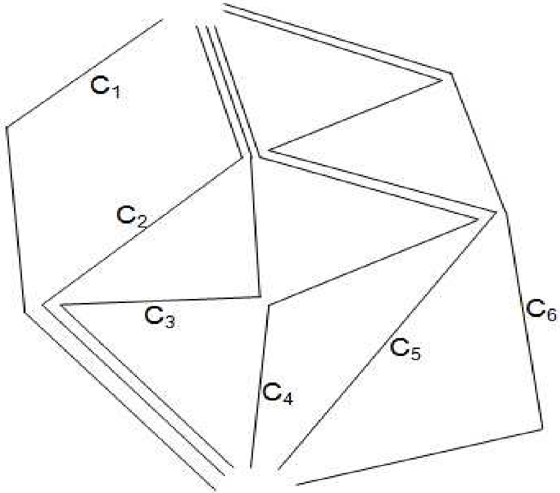
Planar point location by triangulation refinement



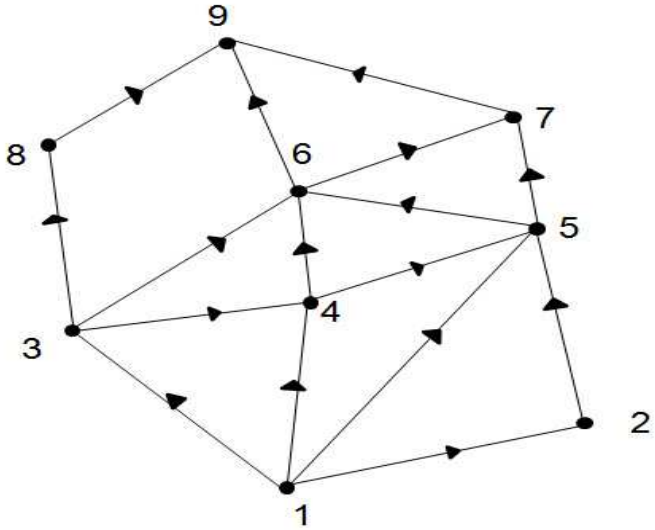
Planar point location using monotone chains



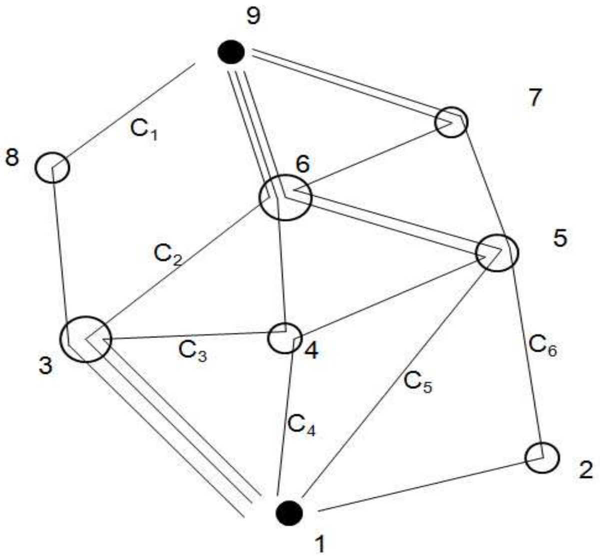
Planar point location using monotone chains



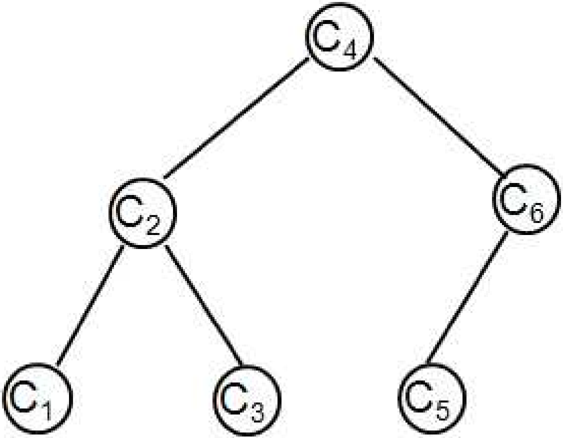
Planar point location using monotone chains



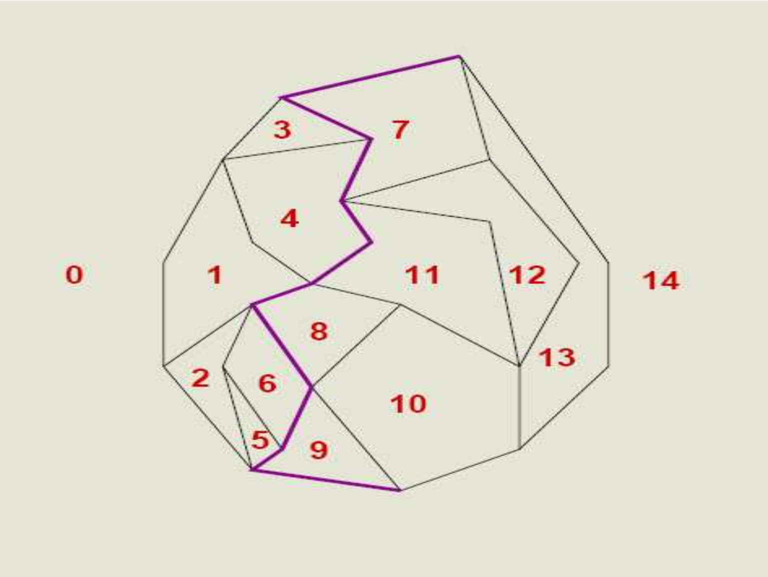
Planar point location using monotone chains



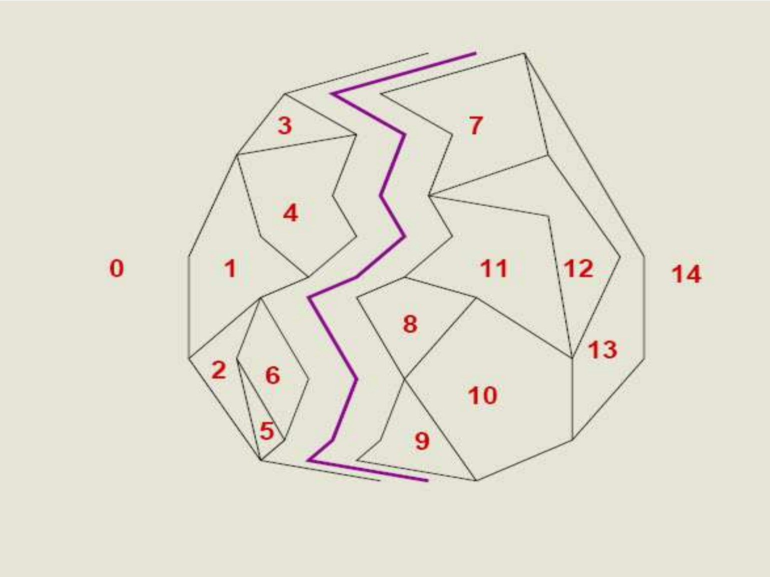
Planar point location using monotone chains



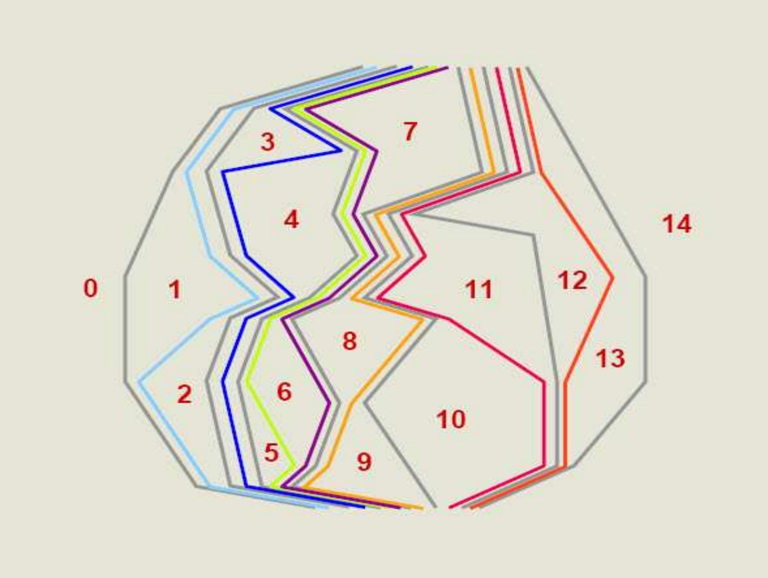
Planar point location using monotone chains



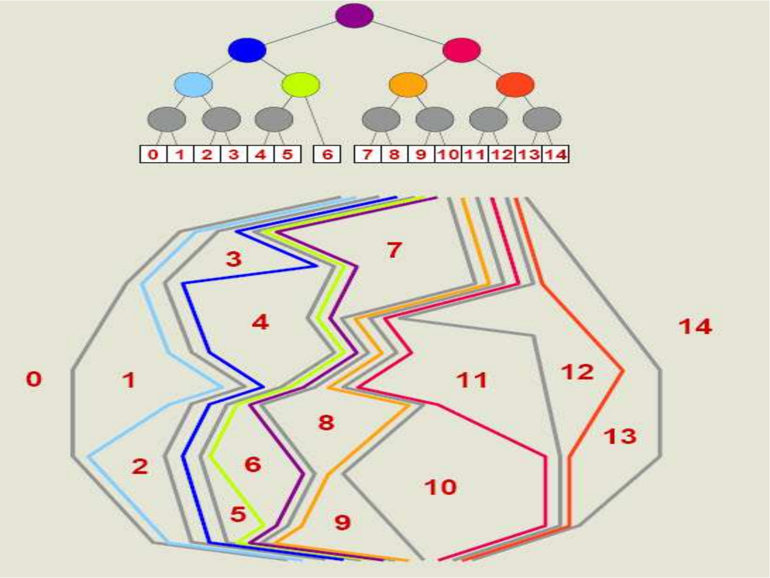
Planar point location using monotone chains



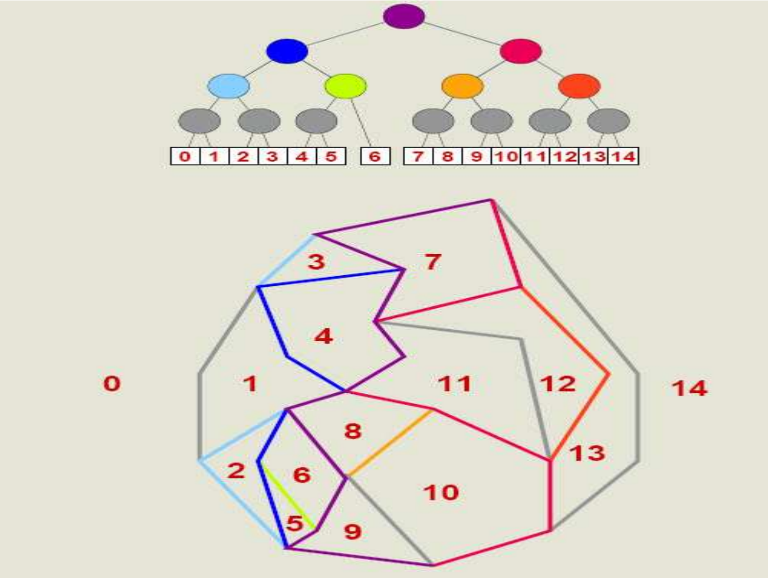
Planar point location using monotone chains



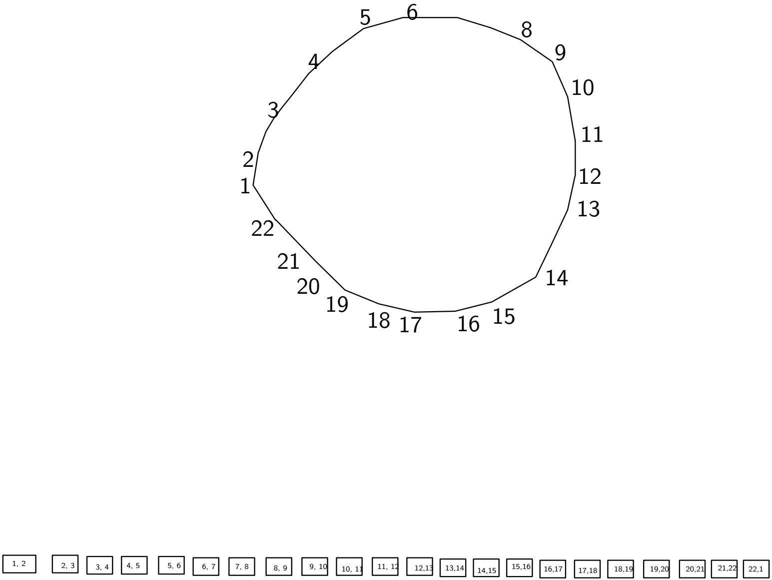
Planar point location using monotone chains



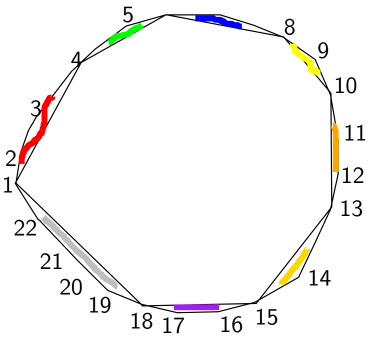
Planar point location using monotone chains

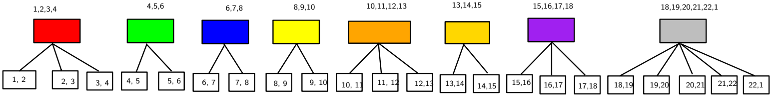


Representing a convex object layer by layer

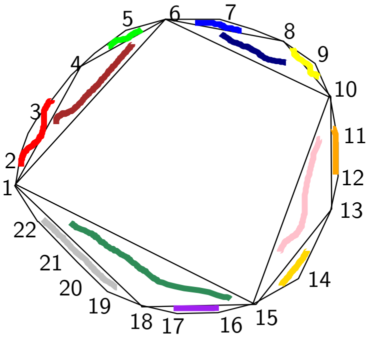


Second layer





Third layer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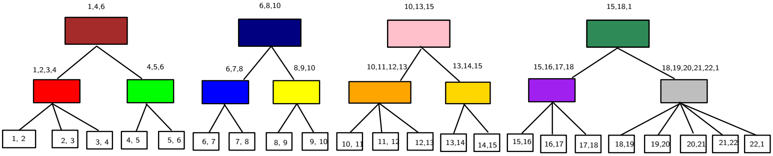
22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10,11,12,13
13,14,15
15,16,17,18
18,19,20,21,22,1
1, 2
2, 3
3, 4
4, 5
5, 6
6, 7
7, 8
8, 9
9, 10
10, 11
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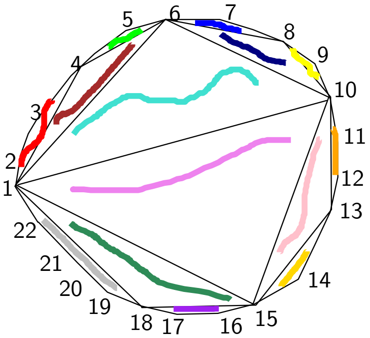
18,19
19,20
20,21
21,22
22,1



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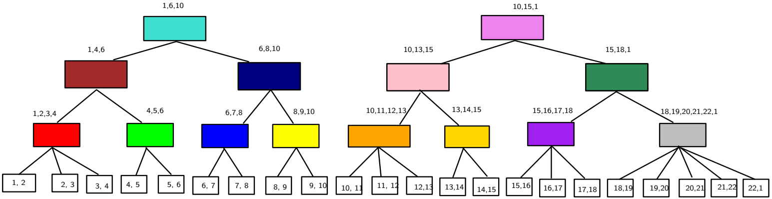

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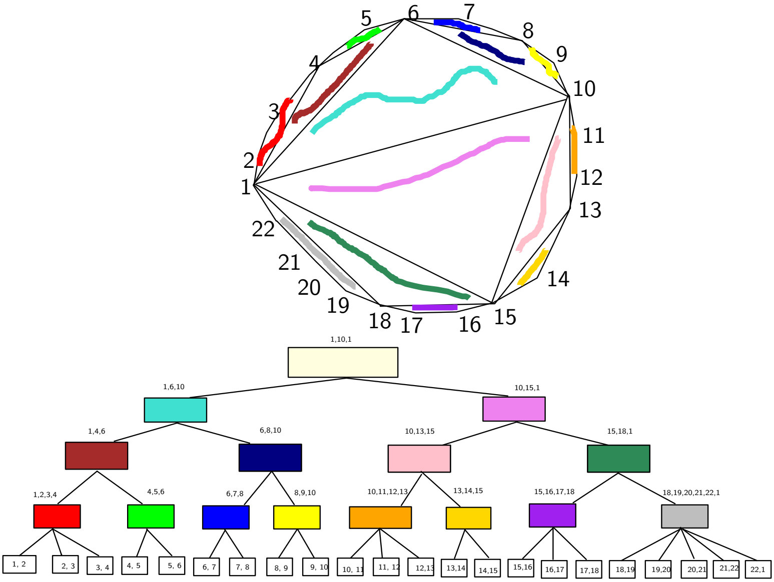
22,1

Point inclusion and Line intersection queries



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